# Improved Spectral Calculations for Discrete Schrödinger Operators 

by

Charles Puelz

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## Approved, Thesis Committee:



Mark Embree, Chair
Professor of Computational and Applied Mathematics,


[^0]Houston, Texas
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# ABSTRACT <br> Improved Spectral Calculations for Discrete Schrödinger Operators 

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This work details an $O\left(n^{2}\right)$ algorithm for computing spectra of discrete Schrödinger operators with periodic potentials. Spectra of these objects enhance our understanding of fundamental aperiodic physical systems and contain rich theoretical structure of interest to the mathematical community. Previous work on the Harper model led to an $O\left(n^{2}\right)$ algorithm relying on properties not satisfied by other aperiodic operators. Physicists working with the Fibonacci Hamiltonian, a popular quasicrystal model, have instead used a problematic dynamical map approach or a sluggish $O\left(n^{3}\right)$ procedure for their calculations. The algorithm presented in this work, a blend of well-established eigenvalue/vector algorithms, provides researchers with a more robust computational tool of general utility. Application to the Fibonacci Hamiltonian in the sparsely studied intermediate coupling regime reveals structure in canonical coverings of the spectrum that will prove useful in motivating conjectures regarding band combinatorics and fractal dimensions.

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## Chapter 1

## Background and Literature Review

Mathematical models of quantum systems present fascinating and subtle computational challenges. This thesis seeks to enable difficult spectral calculations of quantum mechanical operators by proposing an $O\left(n^{2}\right)$ computational tool of general utility and applying this tool to study a mathematical model of a quasicrystal, the Fibonacci Hamiltonian.

A typical finite difference approximation of the time-independent part of the Schrödinger equation, the canonical PDE in quantum mechanics, yields a special type of operator known variously as a Hamiltonian, tight-binding model, or discrete Schrödinger operator. Researchers in different fields study this object in several ways. Physicists interested in numerical simulations of quantum waves restrict the domain of this operator to a finite-dimensional subspace and evolve an initial state on a finite dimensional "lattice" with artificial absorbing boundary conditions. Mathematicians directly handle the infinite-dimensional operator with the goal of proving carefully formulated conjectures. Although different disciplines employ distinct approaches, researchers in both fields agree that spectra of these operators offer crucial insight into the modeled physical systems and contain rich structure interesting in its own right. First, working from the Schrödinger equation, I explain how one arrives at the
quantum model studied in this thesis.

### 1.1 Periodic Schrödinger operators

The Schrödinger Equation

$$
\imath \hbar \frac{\partial}{\partial t} \Psi(x, t)=(\Delta+V(x, t)) \Psi(x, t) .
$$

models the motion of a particle in a quantum mechanical setting. Here $\hbar$ is Planck's constant, $V(x, t)$ is the potential, and $|\Psi(\cdot, t)|^{2}$ is the time dependent probability density of the quantum particle at time $t$. In the following, $\hbar$ is scaled to 1 and the potential is time-independent.

Let $\mathcal{H}:=\Delta+V(x)$. A separation of variables approach with the assumption $\Psi(x, t)=\psi(x) \phi(t)$ results in a constant $E$ and two equations, one in position, and one in time:

$$
\begin{gather*}
\mathcal{H} \psi=E \psi  \tag{1.1}\\
\frac{\partial \phi}{\partial t}=\frac{-\imath E}{\hbar} \phi . \tag{1.2}
\end{gather*}
$$

The first equation is known as the Time Independent Schrödinger Equation (TISE) and the operator $\mathcal{H}$ is called the Hamiltonian. To numerically approximate solutions to (1.1), it is natural to study another operator $\mathbf{H}: \ell^{2}(\mathbb{Z}) \rightarrow \ell^{2}(\mathbb{Z})$, a
discretization of $\mathcal{H}$. Redefine $\psi \in \ell^{2}(\mathbb{Z})$ and take

$$
\begin{equation*}
(\mathbf{H} \psi)_{n}=\psi_{n-1}+v_{n} \psi_{n}+\psi_{n+1} . \tag{1.3}
\end{equation*}
$$

The discrete approximation to the potential is $\left\{v_{n}\right\} \in \ell^{\infty}(\mathbb{Z})$, and a finite difference approximation to the Laplacian $\psi_{n+1}-2 \psi_{n}+\psi_{n-1}$ provides the off-diagonal terms in (1.3). (The scaling of the diagonal by 2 simply shifts the spectrum of $\mathbf{H}$ and is absorbed into the formula for $v_{n}$.) This form of "approximation" to the Hamiltonian, called a Jacobi operator, serves as a fundamental quantum mechanical model and is the central focus of this work.

Definition 1.1. Let $\left\{a_{n}\right\},\left\{b_{n}\right\} \in \ell^{\infty}(\mathbb{Z})$. A Jacobi operator $\mathbf{J}: \ell^{2}(\mathbb{Z}) \rightarrow \ell^{2}(\mathbb{Z})$ is defined as

$$
(\mathbf{J} \psi)_{n}=a_{n-1} \psi_{n-1}+b_{n} \psi_{n}+a_{n} \psi_{n+1}
$$

for all $\psi \in \ell^{2}(\mathbb{Z})$.

With the canonical basis $\left\{\mathbf{e}_{n}\right\}$ defined with 1 in the $n$th component and 0 otherwise, $\mathbf{J}$ can be displayed as the infinite dimensional tridiagonal matrix:

$$
\mathbf{J}=\left[\begin{array}{lllllll}
\ddots & \ddots & \ddots & & & & \\
& a_{-2} & b_{-1} & a_{-1} & & & \\
& & a_{-1} & b_{0} & a_{0} & & \\
& & & a_{0} & b_{1} & a_{1} & \\
& & & & \ddots & \ddots & \ddots
\end{array}\right] .
$$

The sequences $\left\{a_{n}\right\}$ and $\left\{b_{n}\right\}$ are called the Jacobi parameters, and when $\left\{a_{n}\right\}$ is the constant sequence of 1's so $\mathbf{J}$ takes the form of equation (1.3), the operator is called a discrete one-dimensional Schrödinger operator.

A systematic study of the spectrum of particular Schrödinger operators provides significant physical insight into the corresponding quantum systems.

Definition 1.2. The spectrum of $\mathbf{J}$, denoted $\sigma(\mathbf{J})$, is the following set of values:

$$
\sigma(\mathbf{J})=\{\lambda: \mathbf{J}-\lambda \mathbf{I} \text { does not have a bounded, densely defined inverse }\} .
$$

The spectrum of $\mathbf{J}$ can be calculated from finite matrices if the Jacobi parameters satisfy a periodicity condition (see e.g. Teschl [28, pp. 119-121]). My algorithm, described in Chapter 2, comes directly from this reformulation of the spectral calculation as a finite dimensional eigenvalue computation. More precisely, the Jacobi operator $\mathbf{J}$ is called $q$-periodic provided there exists $q \in \mathbb{Z}^{+}$such that $a_{n}=a_{n+q}$ and $b_{n}=b_{n+q}$ for all $n \in \mathbb{Z}$. If such periodicity conditions hold, then the spectrum of $\mathbf{J}$ is the union of $q$ intervals whose endpoints are the eigenvalues of two matrices:

$$
\mathbf{J}_{ \pm}=\left[\begin{array}{ccccc}
b_{1} & a_{1} & & & \pm a_{0} \\
a_{1} & b_{2} & \ddots & & \\
& \ddots & \ddots & \ddots & \\
& & \ddots & b_{q-1} & a_{q-1} \\
\pm a_{0} & & & a_{q-1} & b_{q}
\end{array}\right]
$$

Theorem 1.1. (see e.g. Teschl [28, pp. 119-121]) Let $\mathbf{J}$ be a q-periodic Jacobi operator with parameters $\left\{a_{n}\right\}$ and $\left\{b_{n}\right\}$. Then

$$
\sigma(\mathbf{J})=\bigcup_{k=1}^{q}\left[E_{2 k-1}, E_{2 k}\right]
$$

where the $E_{1}<E_{2} \leq E_{3}<E_{4} \leq \cdots \leq E_{2 q-1}<E_{2 q}$ comprise $\sigma\left(\mathbf{J}_{+}\right) \cup \sigma\left(\mathbf{J}_{-}\right)$.

The quasicrystal model investigated in Chapter 3, formulated as a Jacobi operator, contains aperiodic structure that falls qualitatively between the extremes of periodic and random. In turn, it is relevant to first understand a fundamental distinction between random and periodic quantum systems.

### 1.2 Random and Periodic Quantum Models

The potential sequence $\left\{v_{n}\right\}$ in (1.3) encodes properties of the quantum system to be modeled: researchers have thoroughly investigated the fundamental cases of periodic and random potentials. For an example of the former potential, take $v_{n}=1$ if $n$ is even and 0 if $n$ is odd. For an example of the latter, take $v_{n}$ to be a uniform
random number between 0 and 1 for each $n$. A periodic potential may be interpreted physically as a perfect crystal or pristine physical system, while a random potential is often meant to simulate a system with impurities, or imperfections. Understanding transport through random potentials is significant from a modeling perspective since experimental apparati are often far from ideal systems and contain "random" impurities in their construction. In turn, theoretical models designed to simulate these experiments must take into account these impurities, and a good understanding of their effects in quantum systems is of crucial importance. Not surprisingly, periodic and random models exhibit different electronic transport properties.

The underlying distinction between these two types of models is the following: electrons transport well in periodic systems and poorly in random systems. Transport in periodic systems is called ballistic in certain contexts, and lack of diffusion in random systems is termed Anderson localization after the scientist Phillip Anderson, who wrote the seminal paper in this area [2]. The latter term refers to a literal localization in space of an electron wave spreading through the lattice.

Figure 1.1 displays an initial delta function evolved in time on both a periodic and random lattice. This simulation was performed in Matlab by truncating the Hamiltonian to an order 1001 tridiagonal matrix, and evolving the initial state in time with the matrix exponential as given below:

$$
\Psi(x, t)=\exp \left(-\frac{\imath}{\hbar} \mathbf{H} t\right) \Psi(x, 0)
$$

If the snapshots of the wavefunction on the random lattice are plotted on a loglinear scale, one would observe an exponential envelope encasing the wavefunction. This phenomenon is termed exponential localization, meaning that the tails of the wavefunction decay in a roughly exponential fashion. The snapshots on each row correspond to the same time and were chosen to demonstrate diffusion in the periodic case and, as Anderson described it, the "absence of diffusion" in the random case [2].


Figure 1.1 : Four snapshots in time of an evolving wave. Diffusion occurs in the periodic lattice, and localization occurs in the random lattice.

### 1.3 Quasiperiodic Quantum Models

Several relevant quantum mechanical models have potentials that are determinstic, but not periodic, and so they fall in between the two extremes of periodic and random: these are called quasiperiodic or aperiodic models. It is natural to wonder how electrons move in these aperiodic structures: do they localize like in the random case, diffuse as in the periodic case, or something in between? Answers to this question are often elusive and model dependent. Gaining insight into electronic properties of quasiperiodic systems is the motivation for my computational work discussed in Chapter 2.

Two of the more popular quasiperiodic Schrödinger operator models are the Al most Mathieu operator and the Fibonacci Hamiltonian. The Almost Mathieu operator, also called the Harper model in the physics literature, has the potential

$$
v_{n}=\lambda \cos (2 \pi \sigma n+\theta),
$$

with $\lambda>0$ and $\sigma$ irrational. This potential $\left\{v_{n}\right\}$ is not periodic, and hence provides an example of a sequence in between the periodic and random. In the words of the luminous physicist D.J. Thouless, this model "describe[s] the quantum theory of an electron confined to a plane with a periodic potential in the plane and a uniform magnetic field perpendicular to the plane" [29]. Although not a focus, I mention this model since it has been the center of extensive research, some of which sets a
precedent for my computational work. I will detail this information in the coming sections.

This work focuses on the Fibonacci Hamiltonian, perhaps one of the more popular quasicrystal models. Quasicrystals are peculiar objects that have attracted recent interest in part due to their unique structure: deterministic but lacking translational symmetry. A Penrose tiling of $\mathbb{R}^{2}$, depicted in Figure 1.2, provides a famous example of such a pattern.


Figure 1.2 : Section of a Penrose tiling of the plane. This pattern is an example of a tiling that lacks translational symmetry, like the aperiodic structure of quasicrystals. Picture taken from: http://en.wikipedia.org/wiki/File:Penrose_sun_3.svg and http://en.wikipedia.org/wiki/Penrose_tiling

Nadia Drake has written a nice historical overview of quasicrystals [11]. Synthetic quasicrystals appeared in labs in the 1980s, most notably in the Nobel prize winning work of Dan Shechtman [11. Also during this time, Sütő and others led foundational theoretical research on models of quasicrystals, in particular the Fibonacci Hamiltonian [24, 25]. While experimentalists like Shechtman could produce these objects in labs, a lingering question remained: do quasicrystals "exist in nature" [11]? This question was answered only recently with works published by Steinhardt et al. : these scientists reported their discovery of a rocks found in Russia containing quasicrystals, and concluded they most likely came from an extraterrestrial source [11. This recent discovery, along with foundational experimental and theoretical work begun in the 1980s, has pushed quasicrystals and their corresponding mathematical models to the forefront of research in mathematical physics.

The Fibonacci Hamiltonian is an example of a quasiperiodic model meant to capture the deterministic and nonperiodic structure of quasicrystals. With $\lambda>0$ and $\phi=\frac{1}{2}(1+\sqrt{5})$, the potential is

$$
v_{n}=\lambda \chi_{\left[1-\phi^{-1}, 1\right)}\left(n \phi^{-1} \bmod 1\right)
$$

The coupling constant $\lambda$ scales the strength of the potential relative to the nearestneighbor terms, and much research focuses on the weak coupling $(\lambda \rightarrow 0)$ and strong coupling $(\lambda \rightarrow \infty)$ regimes. Following the notation in the literature, let $\Sigma_{\lambda}$ be the spectrum of the Fibonacci Hamiltonian. As rigorously shown by Sütő, the spectrum
$\Sigma_{\lambda}$ is a Cantor set of measure zero for all $\lambda>0$ [25]. For completeness, I include the following definition (see Willard [33, p. 217]).

Definition 1.3. A Cantor set $\Sigma \subset \mathbb{R}$ is perfect, compact, and totally-disconnected.

Sütő proved several useful facts linking the spectrum of this model to spectra of closely related periodic Schrödinger operators. More precisely, let $\left\{F_{k}\right\}$ be a Fibonacci sequence of numbers $\left(F_{k}=F_{k-1}+F_{k-2}\right)$ with $F_{0}=1$ and $F_{1}=1$, and consider another Schrödinger operator with potential

$$
v_{n}^{(k)}=\lambda \chi_{\left[1-F_{k-1} / F_{k}, 1\right)}\left(n F_{k-1} / F_{k} \bmod 1\right)
$$

In this case, one obtains the periodic Schrödinger operator by replacing the irrational number $\phi$ in the formula for $v_{n}$ with a rational approximation $F_{k} / F_{k-1}$. As explained in the following theorem, this particular choice of a rational approximant yields periodic operators whose spectra, denoted $\sigma_{k}$, relate to the original Cantor spectrum $\Sigma_{\lambda}$ in a well-understood way.

Theorem 1.2. (Sütő 1987, 1989.) For all $\lambda>0, \Sigma_{\lambda}$ is a measure zero Cantor set and the sets $\sigma_{k}$ satisfy

$$
\begin{gather*}
\sigma_{k} \cup \sigma_{k+1} \subset \sigma_{k-1} \cup \sigma_{k}  \tag{1.4}\\
\Sigma_{k}=\bigcap_{k \geq 1} \sigma_{k} \cup \sigma_{k+1} . \tag{1.5}
\end{gather*}
$$

It is known that $\Sigma_{\lambda}$ is a dynamically defined Cantor set for $\lambda$ large enough and
small enough; see Definition 5.2 and Corollary 5.6 in the survey paper by Damanik, Embree, and Gorodetski for a summary of results [8]. Further, the authors conjecture that for all $\lambda>0, \Sigma_{\lambda}$ is dynamically defined. The precise definition given in the survey follows below [8]. Let $B \subset \mathbb{R}$ be a closed interval.

Definition 1.4. (From [8]) A dynamically defined Cantor set $\Sigma \subset B$ satisfies the following:
there exist strictly monotone contracting maps $\psi_{1}, \psi_{2}, \ldots, \psi_{j}: B \rightarrow B$ so that

$$
\psi_{k}(B) \cap \psi_{l}(B)=\emptyset \text { provided } k \neq l,
$$

and with $B_{1}=\psi_{1}(B) \cup \cdots \cup \psi_{j}(B), B_{n+1}=\psi_{1}\left(B_{n}\right) \cup \cdots \cup \psi_{j}\left(B_{n}\right)$ one has

$$
\Sigma=\bigcap_{n \geq 1} B_{n} .
$$

A simple example of such a Cantor set is the famous "middle thirds" Cantor set, obtained by iteratively removing the middle third from a sequence of intervals, and then intersecting the remaining sets. In this case, $B=[0,1]$ and the contracting maps used in the construction are $\psi_{1}(x)=\frac{1}{3} x$ and $\psi_{2}(x)=\frac{1}{3} x+\frac{2}{3}$.

The Cantor structure of the spectrum as formulated in Theorem 1.2 is a profound result and perhaps surprising at first glance. Peering at this operator instead through the lens of dynamical systems enhances one's intuition: values in the spectrum $\Sigma_{\lambda}$ (and covers $\sigma_{k} \cup \sigma_{k+1}$ ) relate to bounded iterates of a particular map. This fact was
also shown by Sütő [24].

Theorem 1.3. (Sütő 1987.)
Fix $\lambda>0$ and $E \in \mathbb{R}$. Let $\left\{x_{k}\right\}$ be a sequence of numbers with the initial conditions $x_{-1}=2, x_{0}=E, x_{1}=E-\lambda$, and satisfying $x_{k+1}=x_{k} x_{k-1}-x_{k-2}$. Then:

$$
\begin{aligned}
& \sigma_{k}=\left\{E:\left|x_{k}(E)\right| \leq 2\right\} \text { and } \\
& \Sigma_{\lambda}=\left\{E: \sup _{k}\left|x_{k}(E)\right|<\infty\right\} .
\end{aligned}
$$

Although the algorithm I present in Chapter 2 is general, the motivation for my computational work applied to the Fibonacci Hamiltonian relies on the theoretical foundation provided by Theorems 1.2 and 1.3 . More precisely, to enable numerical calculations, one takes the covers $\sigma_{k} \cup \sigma_{k+1}$ as approximations to the original Cantor spectrum $\Sigma_{\lambda}$, and then from these covers extracts approximations of interesting quantities. These "interesting quantities," like fractal dimensions, provide physical insight into quasicrystals and thus encourage a careful study of the spectrum $\Sigma_{\lambda}$.

### 1.4 Why study the spectrum?

Physicists and mathematicians glean useful theoretical and physical intuition by studying the spectrum of quantum models. The work of physicists has related fractal dimensions of spectra to time dependent quantities associated with the wave function $\Psi(x, t)$ (see e.g. [17, [18, 23, 32]). Numerical results suggest a connection between
temporal behavior and fractal dimensions of spectra and of the wave functions, although I note that some researchers have provided numerical evidence weakening the former connnection in certain contexts [32]. Regardless, spectral computations are useful in studying "crictical level statisitics" of relevant quantum models [6, 22, 26].

More recently, mathematicians have developed rigorous results for time-dependent behavior: see the work of Damanik et al. for a result incoporating the fractal dimension of the spectrum [9]. In addition, theoreticians work to rigorously explain the mathematical beauty inherent in models and their spectra; the Fibonacci Hamiltonian is a prime example of such a model containing rich theoretical structure, as touched upon in Theorems 1.2 and 1.3. Thus, future work on such models and their spectra will refine our understanding of the electronic properties of quasicrystals and other aperiodic systems, lead to real world applications, and reveal interesting mathematics. My computational work will aide in these endeavors. First, I will discuss previous computational work on quasiperiodic Schrödinger operator models.

### 1.5 Previous computational work on aperiodic models

In this section I detail the computational work done for approximating the spectrum of versions of the Harper model and Fibonacci Hamiltonian. The general theme for numerical calculations is to take a rational approximant to the irrational parameter appearing in the formula for the potential, and then compute the spectrum of the corresponding periodic Schrödinger operator. One hopes that the spectrum of the
periodic operator accurately approximates the original spectrum. More specifically, in the Fibonacci Hamilonian, $\phi$ is approximated by $F_{k} / F_{k-1}$, and in the Harper model, $\sigma$ is approximated by a nearby rational, $p / q$, in both cases resulting in a periodic potential. For the Fibonacci and Harper models, Theorem 1.2 provides a way to tackle the spectral calculation of the periodic operator with an eigenvalue algorithm applied to $\mathbf{J}_{ \pm}$. In the Fibonacci case, Theorem 1.3 provides an additional method to compute the spectrum via a root-finding algorithm for $x_{k}(E)-2$ and $x_{k}(E)+2$ [8].

Obtaining a "good approximation" to the Cantor spectrum (for example, in the sense of equations (1.4) and (1.5) for the Fibonacci Hamiltonian) requires large $k$, or equivalently considering periodic approximations to the potential of increasingly long period. Figures 1.3 and 1.4 demonstrate the convergence of the covers $\sigma_{k} \cup \sigma_{k+1}$ to the Cantor spectrum for the Fibonacci model by displaying exponential decay of interval widths and minimum gap widths. Here, $k=25$, so the largest eigenproblem to be solved is of size $F_{26}=196,418$. Thus, approximating the spectrum $\Sigma_{\lambda}$ via the eigenvalues of $\mathbf{J}_{ \pm}$leads to an extremely large eigenvalue problem. Further, for $\lambda$ sufficiently small, one can can consider even larger problems due to the slower decay of band sizes in the covers $\sigma_{k} \cup \sigma_{k+1}$. Such large-scale computations are not solvable via common algorithms requiring $O\left(n^{3}\right)$ floating point operations (flops). For example, an application of the QR algorithm first requires $O\left(n^{3}\right)$ flops to reduce $\mathbf{J}_{ \pm}$to similar tridiagonal matrices, then $O\left(n^{2}\right)$ flops to compute the eigenvalues.

In light of the challenge of largescale eigenvalue computations, physicists per-


Figure 1.3 : Maximum interval length (top) and minimum interval length (bottom) in $\sigma_{k} \cup \sigma_{k+1}$. Note the exponential decay of the lengths as the covers better approximate the Cantor spectrum $\Sigma_{\lambda}$.


Figure 1.4: Maximum gap length (top) and minimum gap length (bottom) in $\sigma_{k} \cup$ $\sigma_{k+1}$. Note the exponential decay of the minimum gap length as the covers better approximate the Cantor spectum $\Sigma_{\lambda}$. Apparently the maximum gap length remains relatively constant.
forming calculations for Fibonacci-type models have often resorted to a study of latter method: the dynamical map (e.g. [18, 20, 22]). As explained in the paper by Damanik et al., computations in this context present their own challenges [8]. The authors note that in the Fibonacci Hamiltonian model, the polynomial in $E, x_{k}(E)$, has coefficients quickly tending to infinity as $k$ tends to infinity. Further, the slope of $x_{k}(E)$ becomes arbitrarily large in the regions where it is bounded in magnitude by 2 , since the intervals in the spectrum of the periodic approximations converge to a Cantor set. These properties destroy the utility of root-finding calculations.

Researchers have directly tackled the large-scale eigenvalue problem for the Harper model by constructing a similarity transformation converting $\mathbf{J}_{ \pm}$to tridiagonal form; an application of the QR algorithm computes the eigenvalues of this tridiagonal reduction in $O\left(n^{2}\right)$ flops. Partial work toward this end can be seen in the early paper of Thouless [29]. In 1997, Lamoureux published a paper detailing a complete construction of the similarity transformation for a more general class of matrices, and the Harper model fits this mold [21]. Takada et al. presented a solution in 2004, as applied to a different but closely related version of the Harper model, apparently unaware of Lamoureux's work [26]. The work of Lamoureux and Takada et al. rely on special properties of the potential that do not hold for the Fibonacci Hamiltonian.

Since calculation with the dynamical map can be problematic, recent numerical work for the Fibonacci Hamiltonian uses an eigenvalue algorithm applied directly to $\mathbf{J}_{ \pm}$; these computations appear for example in the research of Mandel and Lifshitz and

Damanik, Embree, and Gorodetski [8, 12, 13]. My work provides these researchers and others with a tool to push these numerical computations to higher, more accurate levels (in the sense of approximating the spectrum of the quasiperiodic operator). The algorithm detailed below computes the spectrum of $\mathbf{J}_{ \pm}$in $O\left(n^{2}\right)$ flops with $O(n)$ storage and requires no assumptions on the Jacobi parameters. It can be applied to any periodic Schrödinger operator whose spectrum approximates an aperiodic model. To the best of my knowledge, this is the first available $O\left(n^{2}\right)$ algorithm available for calculation of the covers $\sigma_{k} \cup \sigma_{k+1}$ for the spectrum of the Fibonacci Hamiltonian.

## Chapter 2

## Algorithm for Computing Schrödinger Spectra

dddFollowing from the previous section, to approximate the Cantor spectrum $\Sigma_{\lambda}$ by the canonical covers $\sigma_{k} \cup \sigma_{k+1}$, I wish to compute all the eigenvalues of matrices of the structure

$$
\mathbf{J}_{ \pm}=\left[\begin{array}{ccccc}
b_{1} & a_{1} & & & \pm a_{0}  \tag{2.1}\\
a_{1} & b_{2} & \ddots & & \\
& \ddots & \ddots & \ddots & \\
& & \ddots & b_{q-1} & a_{q-1} \\
\pm a_{0} & & & a_{q-1} & b_{q}
\end{array}\right] .
$$

The proposed algorithm is formulated with the notation given above. Recall that the $2 q$ sorted eigenvalues of (2.1) give the end-points of the intervals comprising the spectrum of the $q$-periodic Jacobi operator with parameters $\left\{a_{n}\right\}$ and $\left\{b_{n}\right\}$. The corners entries $\pm a_{0}$ complicate reduction of $\mathbf{J}_{ \pm}$to tridiagonal form via orthogonal transformations. For example, a naive implementation that uses Givens rotations to "push" the corner entries inward fills in much of the upper and lower portions of the matrix with nonzero values. As noted, working from the structure of the almost Mathieu model, Lamoureux and Takada et al. were able to elegantly avoid this
difficulty by building orthogonal matrices transforming $\mathbf{J}_{ \pm}$to tridiagonal form [21, 26]. These approaches utilize certain properties of the almost-Mathieu potential and are not immediately generalizable to other operators like the Fibonacci Hamiltonian.

Instead of utilizing an explicitly constructed orthogonal transformation, my algorithm splits $\mathbf{J}_{ \pm}$into a tridiagonal plus low-rank modification. This technique maintains generality: it can be used to compute the spectrum of any periodic Jacobi operator, and has a similar flavor to previous work in eigenvalue computations [3, 4, 5, 15].

I will begin with a detailed discussion of the algorithm, called Algorithm 1, followed by some timing results. This chapter concludes with two tests of the algorithm. The first test applies Algorithm 1 to special matrices for which I have a closed form for their eigenvalues, and the second to the Fibonacci Hamiltonian model to compute $\sigma_{k} \cup \sigma_{k+1}$.

### 2.1 Algorithm

This algorithm is presented for the matrix $\mathbf{J}_{+}$; the one for $\mathbf{J}_{-}$is a simple modification.
First, consider the splitting

with $\mathbf{w}=\sqrt{a_{0}}\left(\mathbf{e}_{1}+\mathbf{e}_{n}\right)$.
Let the diagonalization of $\mathbf{T}$ be denoted $\mathbf{Q D Q}{ }^{*}$, with $\mathbf{Q}$ unitary. Then $\mathbf{J}_{+}$is similar to $\mathbf{D}+\mathbf{z z}^{*}$, where $\mathbf{z}=\mathbf{Q}^{*} \mathbf{w}$. With this notation, the algorithm takes the following form. Names of relevant LAPACK routines are included [1].

## Algorithm 1: Computing spectrum of a periodic Jacobi operator

1. Compute all eigenvalues of $\mathbf{T}$ with the QR algorithm (DSTERF.f).
2. Sequentially compute the eigenvectors of $\mathbf{T}$ via inverse iteration, and store the first and last components. The vector $\mathbf{z}$ is calculated from these stored components (DSTEIN.f).
3. For the eigenvalues of $\mathbf{D}$ such that the corresponding component of $\mathbf{z}$ is nonzero, solve the secular equation associated with $\mathbf{D}+\mathbf{z z}^{*}$ with a Newton-type method (DLAED4.f plus code to account for deflation).

### 2.2 Explanation of subparts

It can be shown that each step in Algorithm 1 takes $O\left(n^{2}\right)$ flops. Below, I detail some important facts and references for each subpart: QR iteration, inverse iteration, and root-finding for the secular equation. For details regarding this standard material in numerical linear algebra, see [30, pp. 202-233] and [10].

### 2.2.1 QR algorithm

The QR iteration for eigenvalue computations is an iterative method based on taking the QR factorization of a matrix. At its simplest, the iteration takes the following form.

## QR iteration

1. Reduce $\mathbf{A}$ to upper Hessenberg form $\mathbf{B}$ (tridiagonal if $\mathbf{A}$ is Hermitian).
2. Let $\mathbf{B}^{(0)}:=\mathbf{B}$. Iterate until $\mathbf{B}^{(n)}$ converges to an upper triangular matrix (diagonal if $\mathbf{A}$ is Hermitian):
(a) QR factorization: $\mathbf{Q}^{(n)} \mathbf{R}^{(n)}=\mathbf{B}^{(n-1)}$.
(b) $\mathbf{B}^{(n)}:=\mathbf{R}^{(n)} \mathbf{Q}^{(n)}$. Go to (a).

Several important modifications vastly improve convergence. Implementations of this algorithm factor a shifted matrix at step 2(a), and with minor additions, convergence is achieved in $O\left(n^{2}\right)$ operations. Further if $\lambda_{l}$ and $\tilde{\lambda}_{l}$ are the exact and computed eigenvalues, and $\varepsilon_{\text {mach }}$ is machine epsilon, then

$$
\frac{\left|\lambda_{l}-\tilde{\lambda}_{l}\right|}{\|A\|}=O\left(\varepsilon_{\mathrm{mach}}\right)
$$

I use LAPACK's DSTERF.f, a QR algorithm for symmetric tridiagonal matrices, in my implementation of Algorithm 1.

### 2.2.2 Inverse iteration

The inverse iteration procedure for computing eigenvector components corresponding to an eigenvalue $\gamma$ in second step of Algorithm 1 is simply the power method applied to the matrix $(\mathbf{A}-\alpha \mathbf{I})^{-1}$. Here $\alpha \approx \gamma$, implying the eigenvalue of largest magnitude
is $1 /(\gamma-\alpha)$ and the power method converges to the corresponding eigenvector.

## Inverse iteration

Let $\mathbf{x}^{(0)}$ be an initial normalized vector. Iterate until convergence:

1. Solve $(\mathbf{A}-\alpha \mathbf{I}) \mathbf{y}^{(n)}=\mathbf{x}^{(n-1)}$ for $\mathbf{y}^{(n)}$.
2. $\mathbf{x}^{(n)}=\mathbf{y}^{(n)} /\left\|\mathbf{y}^{(n)}\right\|$. Go to 1 .

Now assume $\mathbf{A}$ is tridiagonal. In this case, steps 1 and $2 \operatorname{cost} O(n)$ flops. With $\alpha$ chosen as the computed eigenvalue from step 1 , the convergence to the corresponding eigenvector is extremely fast, given suitable spacing between the eigenvalues [10, pp. 212, 231]. In turn, the cost for inverse iteration to compute all eigenvectors is $O\left(n^{2}\right)$ flops.

The caveat is the condition of "suitable separation" of the eigenvalues; nearby eigenvalues can yield inaccurately calculated eigenvectors. Further, as depicted in Figures $\sqrt{1.3}$ and 1.4 , the Fibonacci model naturally yields close eigenvalues, since they approximate a set of measure zero. I remark here that I have not experienced difficulties with inverse iteration in the sense of numerical problems in step 2 polluting the final computed eigenvalues, but the proximity of the eigenvalues in my application means we should be vigilant for such problems. Demmel notes that research has indicated the possibility "that inverse iteration may be 'repaired' to provide accurate, orthogonal eigenvectors without spending more than $O(n)$ flops per eigenvector " 10 ,
p. 231]. In turn, future work may result in an improvement of Algorithm 1, step 2, with a more robust but equally efficient algorithm for eigenvectors.

My implementation of Algorithm 1 uses LAPACK's DSTEIN.f, a routine specilized for eigenvector calculations via inverse iteration for symmetric tridiagonal matrices.

### 2.2.3 Secular equation and timing results

Step 3 of Algorithm 1 requires computation of the eigenvalues of a diagonal matrix plus a symmetric rank-one modification. That is, compute $\delta$ so that

$$
\begin{equation*}
\operatorname{det}\left(\mathbf{D}+\mathbf{z z}^{*}-\delta \mathbf{I}\right)=0 \tag{2.2}
\end{equation*}
$$

Several facts allow one to rewrite 2.2 in an equivalent form called the secular equation. Note that in this context, the eigenvalues of the matrices $\mathbf{D}$ and $\mathbf{D}+\mathbf{z z}^{*}$ are real, since the former is similar to a Hermitian matrix and the latter itself is Hermitian.

Lemma 2.1. $\operatorname{det}\left(\mathbf{I}+\mathbf{x y}^{*}\right)=1+\mathbf{y}^{*} \mathbf{x}$. 团

Proof. Begin with the following matrix equation:

$$
\left[\begin{array}{cc}
\mathbf{I} & \mathbf{0} \\
\mathbf{y}^{*} & 1
\end{array}\right]\left[\begin{array}{cc}
\mathbf{I} & -\mathbf{x} \\
\mathbf{0} & 1+\mathbf{y}^{*} \mathbf{x}
\end{array}\right]=\left[\begin{array}{cc}
\mathbf{I} & -\mathbf{x} \\
\mathbf{0} & 1
\end{array}\right]\left[\begin{array}{cc}
\mathbf{I}+\mathrm{xy}^{*} & \mathbf{0} \\
\mathbf{y}^{*} & 1
\end{array}\right]
$$

Taking the determinant of both sides and using the properties

[^1]- $\operatorname{det}(\mathbf{A B})=\operatorname{det}(\mathbf{A}) \operatorname{det}(\mathbf{B})$;
- $\operatorname{det}(\mathbf{A})=$ product of the eigenvalues of $\mathbf{A}$;
results in the equation

$$
1+\mathbf{y}^{*} \mathbf{x}=\operatorname{det}\left[\begin{array}{cc}
\mathbf{I}+\mathbf{x y}^{*} & \mathbf{0} \\
\mathbf{y}^{*} & 1
\end{array}\right]=\operatorname{det}\left(\mathbf{I}+\mathbf{x y}^{*}\right)
$$

Lemma 2.2. Let $d_{1}<d_{2}<\cdots<d_{n}$ and $\delta_{1} \leq \delta_{2} \leq \cdots \leq \delta_{n}$ be the eigenvalues of $\mathbf{D}$ and $\mathbf{D}+\mathbf{z z}^{*}$ respectively, where $\mathbf{D}$ is diagonal. Then, if every component of $\mathbf{z}$ is nonzero, we have $d_{1}<\delta_{1}<d_{2}<\delta_{2}<\cdots<d_{n}<\delta_{n}$.

For a discussion of the latter lemma, see Wilkinson [31, pp. 94-98]. From this point, without loss of generality I assume that all the entries of $\mathbf{z}$ are nonzero. (If the $j$ th component is equal to zero, then $d_{j} \in \sigma\left(\mathbf{D}+\mathbf{z z}^{*}\right)$ with eigenvector $\mathbf{e}_{j}$, and the remaining problem can be split into two independent smaller problems.) This reduction in dimension of the eigenproblem is sometimes called deflation. Let $\delta \in$ $\sigma\left(\mathbf{D}+\mathbf{z z}^{*}\right)$. By Lemma 2.2, $\mathbf{D}-\delta \mathbf{I}$ is invertible and one can write

$$
\mathbf{D}+\mathbf{z z}^{*}-\delta \mathbf{I}=(\mathbf{D}-\delta \mathbf{I})\left(\mathbf{I}+(\mathbf{D}-\delta \mathbf{I})^{-1} \mathbf{z z}^{*}\right)
$$

In turn, by Lemma 2.1, equation 2.2 is satisfied if and only if

$$
\begin{equation*}
g(\delta):=1+\sum_{j=1}^{n} \frac{\left|z_{j}\right|^{2}}{d_{j}-\delta}=0 \tag{2.3}
\end{equation*}
$$

Equation (2.3) is called the secular equation. In this light, the eigenvalues of $\mathbf{D}+\mathbf{z z}^{*}$ can be calculated by a root-finding algorithm applied to this equation. Work towards developing a sufficiently fast root-finder arose from the rank-one modification to the symmetric eigenproblem, and later from the need to find zeros of 2.3 as a substep of the divide-and-conquer eigenvalue algorithm [7]. Thus, Bunch, Nielsen, and Sorensen built an efficient Newton-type iteration by creating a rational function approximant to $g(\delta)$ [5] (summarized in [10, pp. 221-223] and [30, pp. 229-232]). The LAPACK routine DLAED4.f encodes this particular root-finding algorithm and is what I use in my implementation of Algorithm 1.

I supplemented this routine with code to handle the case when a component of $\mathbf{z}$ is zero and the problem deflates. More precisely, if a component $\left|z_{i}\right| \leq \varepsilon_{\text {mach }}$, the code bypasses the root-finding procedure and the corresponding eigenvalue of $\mathbf{D}+\mathbf{z z}^{*}$ is taken to be the unperturbed diagonal entry $d_{i}$ of $\mathbf{D}$.

Timing results for Algorithm 1 applied to compute $\sigma_{k}$ for the Fibonacci Hamiltonian are displayed in Figure 2.1. Interestingly, step 2 and 3 both take greater time than the QR iteration for $\mathbf{T}$. This observation suggests that step 2 and 3 may be implemented in parallel for greater efficiency.


Figure 2.1 : Timing results for Algorithm 1 applied to the Fibonacci model with double precision arithmetic. Displayed are the timing results for individual tests along with the total time.

### 2.3 Testing and some comparisons

### 2.3.1 Gear matrices

In 1969 paper, C.W. Gear carefully calculates closed-form expressions for the eigenvalues of matrices that are tridiagonal plus small modifications to their first and last rows [14]. Such examples include the matrices $\mathbf{J}_{ \pm}$with the Jacobi parameters $a_{n}=0$ and $b_{n}=1$. Using his notation, let an even number $N$ be the order of the matrix.

The eigenvalues are given as follows:

$$
\begin{gathered}
\sigma\left(\mathbf{J}_{+}\right)=\left\{2 \cos \left(\frac{2 k \pi}{N}\right), 1 \leq k<N / 2\right\} \cup\{-2,2\} \\
\sigma\left(\mathbf{J}_{-}\right)=\left\{2 \cos \left(\frac{(2 k-1) \pi}{N}\right), 1 \leq k \leq N / 2\right\}
\end{gathered}
$$

Tables 2.1, 2.2, 2.3, and 2.4 below compare the eigenvalues computed from Algorithm 1 with the theoretical eigenvalues: the first six eigenvalues are displayed along with the error $\dagger$ Note that for these matrices of relatively large order, the eigenvalues computed with Algorithm 1 in double precision match to 15 digits of the theoretical eigenvalues. In the quadruple precision case for $\mathbf{J}_{-}$, the root-finding for the secular equation failed for two eigenvalues: this problem only arose for the Gear example and did not occur in computations for the Fibonacci Hamilonian. I believe this issue is minor, and perhaps some fine tuning in the deflation code would remedy the problem.

[^2]Table 2.1: First several eigenvalues of $\mathbf{J}_{+}$computed with Algorithm 1 double precision, $N=100,000$.

| exact | computed | error |
| :---: | :---: | :---: |
| -2.0000000000000000 | -2.0000000000000004 | 0.0000000000000004 |
| -1.9999999960521582 | -1.9999999960521588 | 0.0000000000000007 |
| - | -1.9999999960521586 | 0.0000000000000004 |
| -1.9999999842086329 | -1.9999999842086331 | 0.0000000000000002 |
| - | -1.9999999842086331 | 0.0000000000000002 |
| -1.9999999644694242 | -1.9999999644694242 | 0.0000000000000000 |
| - | -1.9999999644694242 | 0.0000000000000000 |
| -1.9999999368345323 | -1.9999999368345325 | 0.0000000000000002 |
| - | -1.9999999368345323 | 0.0000000000000000 |
| -1.999999013039569 | -1.9999999013039578 | 0.0000000000000009 |
| - | -1.9999999013039573 | 0.0000000000000004 |

Table 2.2 : First several eigenvalues of $\mathbf{J}_{-}$computed with Algorithm 1 double precision, $N=100,000$.

| exact | computed | error |
| :---: | :---: | :---: |
| -1.9999999990130395 | -1.9999999990130397 | 0.0000000000000002 |
| - | -1.9999999990130395 | 0.0000000000000000 |
| -1.9999999911173560 | -1.9999999911173560 | 0.0000000000000000 |
| - | -1.9999999911173556 | 0.0000000000000004 |
| -1.9999999753259889 | -1.9999999753259894 | 0.0000000000000004 |
| - | -1.9999999753259892 | 0.0000000000000002 |
| -1.9999999516389386 | -1.9999999516389391 | 0.0000000000000004 |
| - | -1.9999999516389384 | 0.0000000000000002 |
| -1.9999999200562049 | -1.9999999200562055 | 0.0000000000000007 |
| - | -1.9999999200562051 | 0.0000000000000002 |
| -1.9999998805777879 | -1.9999998805777885 | 0.0000000000000007 |
| - | -1.9999998805777881 | 0.0000000000000002 |

Table 2.3 : First several eigenvalues of $\mathbf{J}_{+}$computed with Algorithm 1 in quadruple precision, $N=100,000$.

| exact | computed | error |
| :---: | :---: | :---: |
| -2.0000000000000000000000000000000000 | -2.0000000000000000000000000000000000 | $0.00 E+00$ |
| -1.9999999960521582408630444327486558 | -1.9999999960521582408630444327486560 | $0.20 E-33$ |
| - | -1.9999999960521582408630444327486556 | $0.20 E-33$ |
| -1.9999999842086329790376322861801951 | -1.9999999842086329790376322861801953 | $0.20 E-33$ |
| - | -1.9999999842086329790376322861801949 | $0.20 E-33$ |
| -1.9999999644694242612801271643224253 | -1.9999999644694242612801271643224254 | $0.20 E-33$ |
| - | -1.9999999644694242612801271643224253 | $0.00 E+00$ |
| -1.9999999368345321655178015354586642 | -1.9999999368345321655178015354586653 | $0.12 E-32$ |
| - | -1.9999999368345321655178015354586644 | $0.20 E-33$ |
| -1.9999999013039568008488364244832007 | -1.9999999013039568008488364244832012 | $0.60 E-33$ |
| - | -1.9999999013039568008488364244832009 | $0.20 E-33$ |

Table 2.4 : First several eigenvalues of $\mathbf{J}_{-}$computed with Algorithm 1 in quadruple precision, $N=100,000$.

| exact | computed | error |
| :---: | :---: | :---: |
| -1.9999999990130395599722383806422158 | -1.9999999990130395599722383806422161 | $0.40 E-33$ |
| - | -1.9999999990130395599722383806422159 | $0.20 E-33$ |
| -1.9999999911173560455946908858973095 | -1.9999999911173560455946908858973103 | $0.80 E-33$ |
| - | -1.9999999911173560455946908858973097 | $0.20 E-33$ |
| -1.9999999753259890480105049913964137 | -1.9999999753259890480105049913964142 | $0.60 E-33$ |
| - | -1.9999999753259890480105049913964141 | $0.40 E-33$ |
| -1.9999999516389386295614987640595453 | -1.9999999516389386295614987640595461 | $0.80 E-33$ |
| - | -1.9999999516389386295614987640595457 | $0.40 E-33$ |
| -1.9999999200562048837603989966221883 | -1.9999999200562048837603989966221887 | $0.40 E-33$ |
| - | -1.9999999200562048837603989966221883 | $0.00 E+00$ |
| -1.9999998805777879352908408384618465 | -1.9999998805777879352908408384618467 | $0.20 E-33$ |
| - | -1.9999998805777879352908408384618465 | $0.00 E+00$ |

### 2.3.2 Fibonacci model

I now describe numerical results for the Fibonacci Hamiltonian. Most of the calculations in this thesis fall in the less-studied intermediate coupling regime, where $\lambda$ is neither close to zero nor very large (i.e. $\lambda$ is bounded away from 0 but less than 4). This range of $\lambda$ values allows for implementation of Algorithm 1 in double precision, since band widths shrink more slowly for smaller values of $\lambda$ (see Figures 1.3 and 1.4). Thus, calculations may be taken to higher levels than shown in this work, as long as $\lambda$ is sufficiently small to avoid exponentially close eigenvalues. In turn, quadruple precision calculations may be used to tackle larger values of the coupling strength (greater than 4) to resolve differences in close-by eigenvalues. It is possible to compile LAPACK routines to run in quadruple precision, but computational time increases in this extended precision.

I have included Tables 2.5, 2.6, 2.7, and 2.8 displaying the 5 middle computed eigenvalues of the matrices $\mathbf{J}_{ \pm}$for the Fibonacci Hamiltonian with $k=20$ and $\lambda=$ 2,8 . Values in the middle of the spectrum are chosen given the qualitative assumption that substantial clustering occurs in this region. Four different methods are shown: Algorithm 1 in double and quadruple precision ("Alg. 1d" and "Alg. 1q" respectively in the tables below), MATLAB's eig, and the LAPACK routine DSYEV.f compiled to run in quadruple precision.

Note that the eigenvalues for this application will be close together, since they approximate the Cantor spectrum of measure zero of the Fibonacci Hamiltonian. One

Table 2.5 : Five middle eigenvalues in $\sigma\left(\mathbf{J}_{+}\right)$for the Fibonacci Hamiltonian with $\lambda=2$. The order of $\mathbf{J}_{+}$is $10,946(k=20)$. Displayed are the computed eigenvalues from four separate codes, two in double precision and two in quadruple precision.

| Alg. 1d | 1.4292357574320784 |
| :---: | :---: |
| eig | 1.429235757432107 |
| Alg. 1q | 1.4292357574320782330720391457131554 |
| DSYEV.f | 1.4292357574320782330720391457131529 |
| Alg. 1d | 1.4292533308961350 |
| eig | 1.429253330896162 |
| Alg. 1q | 1.4292533308961349521011449450318725 |
| DSYEV.f | 1.4292533308961349521011449450318734 |
| Alg. 1d | 1.4292880261998258 |
| eig | 1.429288026199844 |
| Alg. 1q | 1.4292880261998256521516529349150777 |
| DSYEV.f | 1.4292880261998256521516529349150796 |
| Alg. 1d | 1.4293062009721220 |
| eig | 1.429306200972142 |
| Alg. 1q | 1.4293062009721217677616261343470380 |
| DSYEV.f | 1.4293062009721217677616261343470403 |
| Alg. 1d | 1.4293408978789590 |
| eig | 1.429340897878983 |
| Alg. $1 q$ | 1.4293408978789593473717581090306448 |
| DSYEV.f | 1.4293408978789593473717581090306467 |

obtains more digits of accuracy for $\lambda=2$ as compared to $\lambda=8$, since the eigenvalue spacing decays more gradually with the level number $k$ for smaller values of $\lambda$. The digits highlighted in grey are those of the eigenvalues computed in double precision that agree with the digits of the eigenvalues computed in quadruple precision.

Table 2.6 : Five middle eigenvalues in $\sigma\left(\mathbf{J}_{+}\right)$for the Fibonacci Hamiltonian with $\lambda=8$. The order of $\mathbf{J}_{+}$is $10,946(k=20)$. Notice that these eigenvalues are closer together than for $\lambda=2$ in Table 2.5 .

| Alg. 1d | 7.1233218906775511 |
| ---: | :--- |
| eig | 7.123321890679813 |
| Alg. 1q | 7.1233218906775077259052027990441278 |
| DSYEV. | 7.1233218906775077259052027990441371 |
| Alg. 1d | 7.1233218908967313 |
| eig | 7.123321890898212 |
| Alg. 1q | 7.1233218908966885297296485295470743 |
| DSYEV.f | 7.1233218908966885297296485295470851 |
| Alg. 1d | 7.1233218938161196 |
| eig | 7.123321893816916 |
| Alg. 1q | 7.1233218938160852861910276823143882 |
| DSYEV.f | 7.1233218938160852861910276823144151 |
| Alg. 1d | 7.1233218939197620 |
| eig | 7.123321893920583 |
| Alg. 1q | 7.1233218939197272777658699006219234 |
| DSYEV.f | 7.1233218939197272777658699006219527 |
| Alg. 1d | 7.1233218968391663 |
| eig | 7.123321896840594 |
| Alg. 1q | 7.1233218968391240547230659524326977 |
| DSYEV.f | 7.1233218968391240547230659524327108 |

Table 2.7: Five middle eigenvalues in $\sigma\left(\mathbf{J}_{-}\right)$for the Fibonacci Hamiltonian with $\lambda=2$. The order of $\mathbf{J}_{-}$is $10,946(k=20)$.

| Alg. 1d | 1.4292404547761508 |
| :---: | :---: |
| eig | 1.429240454776190 |
| Alg. 1q | 1.4292404547761508434229558244154517 |
| DSYEV.f | 1.4292404547761508434229558244154530 |
| Alg. 1d | 1.4292493242116628 |
| eig | 1.429249324211683 |
| Alg. 1q | 1.4292493242116629518885897087185001 |
| DSYEV.f | 1.4292493242116629518885897087185007 |
| Alg. 1d | 1.4292934828231283 |
| eig | 1.429293482823148 |
| Alg. 1q | 1.4292934828231283292165371593409164 |
| DSYEV.f | 1.4292934828231283292165371593409164 |
| Alg. 1d | 1.4293007443524455 |
| eig | 1.429300744352459 |
| Alg. 1q | 1.4293007443524456179224176321016020 |
| DSYEV.f | 1.4293007443524456179224176321016049 |
| Alg. 1d | 1.4293449047206941 |
| eig | 1.429344904720721 |
| Alg. 1q | 1.4293449047206944219368822628772715 |
| DSYEV.f | 1.4293449047206944219368822628772715 |

Table 2.8: Five middle eigenvalues in $\sigma\left(\mathbf{J}_{-}\right)$for the Fibonacci Hamiltonian with $\lambda=8$. The order of $\mathbf{J}_{-}$is $10,946(k=20)$.

| Alg. 1d | 7.1233218906885654 |
| :---: | :---: |
| eig | 7.123321890690949 |
| Alg. 1q | 7.1233218906885593314445002636774535 |
| DSYEV.f | 7.1233218906885593314445002636773942 |
| Alg. 1d | 7.1233218908903284 |
| eig | 7.123321890891762 |
| Alg. 1q | 7.1233218908903082601007337396589646 |
| DSYEV.f | 7.1233218908903082601007337396589792 |
| Alg. 1d | 7.1233218938274385 |
| eig | 7.123321893828312 |
| Alg. $1 q$ | 7.1233218938274520441700214641767623 |
| DSYEV.f | 7.1233218938274520441700214641767823 |
| Alg. 1d | 7.1233218939083329 |
| eig | 7.123321893909189 |
| Alg. 1q | 7.1233218939083605197868768550743002 |
| DSYEV.f | 7.1233218939083605197868768550743202 |
| Alg. 1d | 7.1233218968455265 |
| eig | 7.123321896847012 |
| Alg. 1q | 7.1233218968455043244062521287398769 |
| DSYEV.f | 7.1233218968455043244062521287398831 |

## Chapter 3

## Application: Fibonacci Hamiltonian

This chapter details an application of Algorithm 1 for computing the canonical covers $\sigma_{k} \cup \sigma_{k+1}$ of the spectrum $\Sigma_{\lambda}$ of the Fibonacci Hamiltonian. These calculations allow for extrapolating approximations to fractal dimensions from the canonical covers and formulating combinatorial statements for the bands in $\sigma_{k}$ in the coupling regime $\lambda \in(0,4]$. Approximate dimensions inform conjectures regarding the regularity and mononoticity of the exact dimensions as functions of $\lambda$. Combinatorial results aid in approximating dimensions from the canonical covers, increase our physical understanding of quasicrystals, and unveil beautiful structure of theoretical interest. Relevant examples include the work of Killip et al., who present a complete description of the combinatorics for $\lambda>4$ and then apply their result to derive bounds on quantum particle dynamics [19]. Similarly, Damanik, Embree, Gorodetski and Tcheremchantsev utilize these combinatorics to extrapolate bounds on fractal dimensions as $\lambda \rightarrow \infty$; they then relate the box-counting dimension to the time evolution of an initial delta function [9]. My numerical work arises from the same motivations, but focuses on the less-studied intermediate coupling regime. Section 3.1 contains work on fractal dimensions and section 3.2 examines interval combinatorics.

### 3.1 Fractal dimensions

There are several ways to measure the dimension of a fractal subset of $\mathbb{R}$. I study the box-counting and Hausdorff dimensions. Let $A \subset \mathbb{R}$.

Definition 3.1. Let $\varepsilon>0$ and $M_{A}(\varepsilon)=\#\{n \in \mathbb{Z}:[n \varepsilon,(n+1) \varepsilon) \cap A \neq \emptyset\}$. The upper and lower box-counting dimensions of $A$ are respectively

$$
\begin{aligned}
& \operatorname{dim}_{B}^{+} A:=\limsup _{\varepsilon \rightarrow 0} \frac{\log M(\varepsilon)}{-\log \varepsilon} \\
& \operatorname{dim}_{B}^{-} A:=\liminf _{\varepsilon \rightarrow 0} \frac{\log M(\varepsilon)}{-\log \varepsilon} .
\end{aligned}
$$

When the two limits above are equal, their limit, denoted $\operatorname{dim}_{B} A$, is the box-counting dimension of $A$.

Definition 3.2. Let $\Delta>0$ and consider a set of intervals $\left\{C_{l}\right\}_{l \geq 1}$, such that $A \subset$ $\bigcup_{l \geq 1} C_{l}$ and $\left|C_{l}\right| \leq \Delta$. Given

$$
h^{\alpha}(A):=\lim _{\Delta \rightarrow 0} \inf _{\Delta \text {-covers }} \sum_{l \geq 1}\left|C_{l}\right|^{\alpha}
$$

with $\alpha \in[0,1]$, the Hausdorff dimension is

$$
\operatorname{dim}_{H} A:=\inf \left\{\alpha: h^{\alpha}(A)<\infty\right\}
$$

The fractal dimensions of the spectrum $\Sigma_{\lambda}$ reveal electronic transport properties of the Fibonacci Hamiltonian. For example, Theorem 3 in the paper of Damanik
et al. describes bounds the evolution of a portion of a wavepacket by a quantity involving a fractal dimension [9]. In addition, these dimensions give insight into the topological characteristics of spectra of two and three dimensional Fibonacci models. More precisely, higher dimensional models have spectra that are set sums of the onedimensional spectrum $\Sigma_{\lambda}$. For two and three dimensions,

$$
\begin{aligned}
& \sigma\left(\mathbf{H}_{2 d, \lambda}\right)=\Sigma_{\lambda}+\Sigma_{\lambda} \\
& \sigma\left(\mathbf{H}_{3 d, \lambda}\right)=\Sigma_{\lambda}+\Sigma_{\lambda}+\Sigma_{\lambda} .
\end{aligned}
$$

Since the one dimensional spectrum is a Cantor set, what can be said about the set sum of the spectrum with itself? Damanik and Gorodetski relate the structure of the spectrum of the two-dimensional model with the box-counting dimension: $\sigma\left(\mathbf{H}_{2 d, \lambda}\right)$ is a Cantor set provided the upper box-counting dimension is less than $1 / 2$ [8]. Further, more subtle scenarios arise: the sum of a Cantor set with itself may form a Cantorval, an object with similarities to Cantor set and intervals. Definitions and related open problems may be found in the last section of the survey article [8].

Lastly, questions regarding regularity (and mononiticity) of the dimensions (with respect to $\lambda$ ) and the relationship between the two dimensions in the intermediate coupling regime remain open. In this section I discuss approximating dimensions of $\Sigma_{\lambda}$ from the covers $\sigma_{k} \cup \sigma_{k+1}$ with a goal of shedding light on these open problems.

### 3.1.1 Approximating the Hausdorff dimension

I develop approximations to the Hausdorff dimension from the covers $\sigma_{k} \cup \sigma_{k+1}$ in a way derived from the work of Halsey et al. [16]. Papers from Kohmoto et al. and Tang et al. demonstrate an application of this approach to quasiperiodic models [20, 27], although they compute a different but related object (the " $f-\alpha$ " curve of the spectrum).

To approximate the Hausdorff dimension, let $\left\{B_{j, k}\right\}_{j=1}^{N_{k}}$ be an enumeration of the bands in the cover $\sigma_{k} \cup \sigma_{k+1}$. With the bands $\sigma_{k} \cup \sigma_{k+1}$ and $\sigma_{k+1} \cup \sigma_{k+2}$, construct the following function:

$$
g_{k}(\alpha)=\sum_{j=1}^{N_{k}}\left|B_{j, k}\right|^{\alpha}-\sum_{j=1}^{N_{k+1}}\left|B_{j, k+1}\right|^{\alpha} .
$$

The approximate Hausdorff dimension is $\tilde{\alpha}_{k}$ such that $g_{k}\left(\tilde{\alpha}_{k}\right)=0$. One can view $\sigma_{k} \cup \sigma_{k+1}$ and $\sigma_{k+1} \cup \sigma_{k+2}$ as $\Delta$-covers, where $k$ is taken large enough so that $\Delta$ is sufficiently small. In turn, the assumption that both sums "suitably approximate" the limit $h^{\alpha}\left(\Sigma_{\lambda}\right)$ implies:

$$
\frac{\sum_{j=1}^{N_{k}}\left|B_{j, k}\right|^{\alpha}}{\sum_{j=1}^{N_{k+1}}\left|B_{j, k+1}\right|^{\alpha}} \approx 1
$$

Enforcing equality of this ratio, I seek an approximate Hausdorff dimension as a root of $g_{k}$ in $[0,1]$. This root-finding procedure is done with Matlab's fzero function. Table 3.1 depicts the approximate dimension calculated from successively higher levels for four values of $\lambda$; one observes convergence as $k$ increases since the sets $\sigma_{k} \cup \sigma_{k+1}$

Table 3.1: Approximate Hausdorff dimension $\tilde{\alpha}_{k}$ computed with the covers $\sigma_{k} \cup \sigma_{k+1}$ and $\sigma_{k+1} \cup \sigma_{k+2}$. Note the convergence of the approximation to the dimension as the level $k$ increases.

|  | $\lambda=1$ | $\lambda=2$ | $\lambda=3$ | $\lambda=4$ |
| :---: | :---: | :---: | :---: | :---: |
| 3 | 0.6291847216 | 0.5462641412 | 0.5509223280 | 0.4880251151 |
| 4 | 0.7717876114 | 0.6387171661 | 0.4962734864 | 0.4360752232 |
| 5 | 0.7687008005 | 0.5980307042 | 0.5245905225 | 0.4643367675 |
| 6 | 0.7853861811 | 0.6175207019 | 0.5099061560 | 0.4489717534 |
| 7 | 0.7707783314 | 0.6081066010 | 0.5177173248 | 0.4575388294 |
| 8 | 0.7803899153 | 0.6128945234 | 0.5135325765 | 0.4527265647 |
| 9 | 0.7767175230 | 0.6104874494 | 0.5157963293 | 0.4554547177 |
| 10 | 0.7714623010 | 0.6117021346 | 0.5145673349 | 0.4539025279 |
| 11 | 0.7735535776 | 0.6110879386 | 0.5152374261 | 0.4547889870 |
| 12 | 0.7725063848 | 0.6114005132 | 0.5148714885 | 0.4542820205 |
| 13 | 0.7729999110 | 0.6112412695 | 0.5150716978 | 0.4545724092 |
| 14 | 0.7728023687 | 0.6113225062 | 0.5149620869 | 0.4544059893 |
| 15 | 0.7728869425 | 0.6112810399 | 0.5150221437 | 0.4545014247 |
| 16 | 0.7728465217 | 0.6113022263 | 0.5149892285 | 0.4544466855 |
| 17 | 0.7728647751 | 0.6112913969 | 0.5150072741 | 0.4544780900 |
| 18 | 0.7728569242 | 0.6112969343 | 0.5149973793 | 0.4544600723 |
| 19 | 0.7728604870 | 0.6112941023 | 0.5150028057 | 0.4544704110 |
| 20 | 0.7728588318 | 0.6112955508 | 0.5149998295 | 0.4544644829 |
| 21 | 0.7728595719 | 0.6112948098 | 0.5150014635 | 0.4544678747 |
| 22 | 0.7728592431 | 0.6112951888 | 0.5150005739 | 0.4544659127 |
| 23 | 0.7728593933 | 0.6112949949 | 0.5150010568 | 0.4544670354 |
| 24 | 0.7728593247 | 0.6112950942 | 0.5150007784 | 0.4544653981 |

better approximate the Cantor spectrum.
Figure 3.1 shows the approximate Hausdorff dimension as a function of $\lambda$. At this level of resolution in $\lambda$, the dimension appears to be monotone and contains no jumps; these observations support conjectures regarding regularity and monotonicity (i.e. Corollary 5.6 of [8] for a statement about regularity).

Damanik et al. derived upper and lower bounds on the the Hausdorff dimension for $\lambda>16$ [9]. In their notation, with $S_{u}(\lambda)=\frac{1}{2}\left(\lambda-4+\sqrt{(\lambda-4)^{2}-12}\right)$ and


Figure 3.1: Approximate Hausdorff dimension calculated from the covers $\sigma_{k} \cup \sigma_{k+1}$ with $k=23,24$. At this resolution in $\lambda$, the dimension appears smooth and monotone.


Figure 3.2: Approximate Hausdorff dimension calculated from the covers $\sigma_{k} \cup \sigma_{k+1}$ with $k=13,14$. Also plotted are the upper and lower bounds proven in [9].
$S_{l}(\lambda)=2 \lambda+22$, they proved

$$
\operatorname{dim}_{H} \Sigma_{\lambda} \geq \frac{\log (1+\sqrt{2})}{\log S_{u}(\lambda)}
$$

and

$$
\operatorname{dim}_{H} \Sigma_{\lambda} \leq \frac{\log (1+\sqrt{2})}{\log S_{l}(\lambda)}
$$

As a check of this method of approximation, I have plotted the approximate dimension and the bounds in Figure 3.2.

### 3.1.2 Box-counting dimension

Several conjectures regarding the box-counting dimension of $\Sigma_{\lambda}$ require some numerical insight. For both $\lambda \rightarrow 0$ and $\lambda>16$, the box-counting and Hausdorff dimensions are equal, but less is known in the intermediate coupling regime (see [8]). Generally, the fractal dimensions of a set $A \subset \mathbb{R}$ satisfy

$$
\operatorname{dim}_{H} A \leq \operatorname{dim}_{B}^{-} A \leq \operatorname{dim}_{B}^{+} A .
$$

For the spectrum $\Sigma_{\lambda}$, Damanik et al. conjecture the equality of the dimensions for any $\lambda>0$ [8].

In addition, numerical work may reveal critical $\lambda$ values at which transitions in the spectrum of higher dimensional models occur, i.e. $\lambda$ for which the upper boxcounting dimension is equal to $1 / 2$. Assuming $\operatorname{dim}_{B} \Sigma_{\lambda}=\operatorname{dim}_{H} \Sigma_{\lambda}$ for all $\lambda$, Figure

Table 3.2: Approximation of $\lambda$ such that $\operatorname{dim} \Sigma_{\lambda}=1 / 2$. This critical coupling strength is determined by approximating the Hausdorff dimension as discussed in section 3.1.1, and then determining a root of $\operatorname{dim}_{H} \Sigma_{\lambda}-1 / 2$ in the interval $[3,4]$.

| $k$ | root of $\operatorname{dim}_{H} \Sigma_{\lambda}-1 / 2$ |
| :---: | :---: |
| 5 | 3.358132165787465 |
| 6 | 3.136023908556593 |
| 7 | 3.252443323604985 |
| 8 | 3.188369245375197 |
| 9 | 3.223061231166874 |
| 10 | 3.203897752346325 |
| 11 | 3.214420790092100 |
| 12 | 3.208597301692052 |
| 13 | 3.211813603439293 |
| 14 | 3.210032038473626 |
| 15 | 3.211018258537230 |
| 16 | 3.210471722049607 |
| 17 | 3.210774542902178 |
| 18 | 3.210606689833002 |
| 19 | 3.210699725856641 |
| 20 | 3.210648145088315 |
| 21 | 3.210676725662388 |
| 22 | 3.210661023369781 |

3.1 suggests the dimension takes the value $1 / 2$ for $\lambda \approx 3$. Table 3.2 more precisely investigates this critical value of $\lambda$ in the following way. I created a function taking $\lambda \mapsto\left(\right.$ approximation of $\left.\operatorname{dim}_{H} \Sigma_{\lambda}\right)$ by efficiently computing the covers $\sigma_{k} \cup \sigma_{k+1}$ and $\sigma_{k+1} \cup \sigma_{k+2}$ with Algorithm 1, and then from these covers, extracting an approximation to the Hausdorff dimension as described in section 3.1.1. The values displayed in Table 3.2 are the roots of this function in the interval [3, 4].

The relevant data collected from the covers $\sigma_{k} \cup \sigma_{k+1}$ for approximating the boxcounting dimension are particular sequences of infima for each $\lambda$. More precisely,
define the function (refer to Definition 3.1)

$$
f_{k, \lambda}(\varepsilon)=\frac{\log M_{\sigma_{k} \cup \sigma_{k+1}}(\varepsilon)}{-\log \varepsilon} .
$$

From Figures 3.3 and 3.4 , one is hopeful that the sequence $\inf _{\varepsilon \in(0,1)} f_{k, \lambda}(\varepsilon)$ converges to the box-counting dimension of $\Sigma_{\lambda}$. These figures appear in the survey paper by Damanik, Embree, and Gorodetski, but without averaging over $\varepsilon$-grids. Their plots of $f_{k, \lambda}$ appear highly discontinuous, especially for larger values of $\epsilon$. I compute the value of the function $f_{k, \lambda}$ at $\varepsilon$ with averaging by shifting the $\varepsilon$-grid by $(j-1) \varepsilon / N$, where $N$ is the number of realizations. For each $j=1 \ldots N$, I compute a number $-\log M_{j} / \log \varepsilon$, where $M_{j}$ is the number of intervals in the shifted grid that nontrivially intersect $\sigma_{k} \cup \sigma_{k+1} ; f_{k, \lambda}(\varepsilon)$ is taken as the average over these values. With averaging, the resulting plots of $f_{k, \lambda}$ appear smoother and thus facilate more accurate calculation of the $\inf _{\varepsilon \in(0,1)} f_{k, \lambda}(\varepsilon)$.

The approximate Hausdorff dimension as computed in section 3.1.1 is plotted as a dashed line in Figures 3.3 and 3.4 with the plots of $f_{k, \lambda}$ to visualize the connection between the Hausdorff and box-counting dimensions. In all cases the dashed line provides a lower bound on the infima. Table 3.3 displays the sequence of minima for four values of the coupling strength: the convergence is agonizingly slow. From Figure 3.4 one expects quicker convergence for the largest value $\lambda=4$, but even in this case I think we cannot expect even one digit of accuracy (for example, the first digit should be 4 if we assume equality of Hausdorff and box-counting). I have

Table 3.3 : Minima of $f_{k, \lambda}$. Observe the slow convergence.

| $k$ | $\lambda=1$ | $\lambda=2$ | $\lambda=3$ | $\lambda=4$ |
| :---: | :---: | :---: | :---: | :---: |
| 3 | 1.0336903744 | 1.0209104937 | 1.0076755514 | 0.9743044183 |
| 4 | 1.0298323720 | 1.0121711018 | 0.9750578638 | 0.8869700670 |
| 5 | 1.0257606630 | 1.0033435987 | 0.9105012908 | 0.8142126618 |
| 6 | 1.0217885452 | 0.9773356592 | 0.8530686632 | 0.7561537168 |
| 7 | 1.0177107027 | 0.9389940411 | 0.8076428267 | 0.7123968591 |
| 8 | 1.0136620162 | 0.9051927594 | 0.7713904183 | 0.6781950504 |
| 9 | 1.0096087392 | 0.8762657167 | 0.7424399454 | 0.6516463581 |
| 10 | 1.0055581702 | 0.8516590139 | 0.7191551380 | 0.6306586165 |
| 11 | 1.0015053675 | 0.8307835418 | 0.6997014118 | 0.6132573124 |
| 12 | 0.9926301691 | 0.8127909835 | 0.6836084473 | 0.5990907457 |
| 13 | 0.9801915271 | 0.7972557922 | 0.6698542777 | 0.5871119859 |
| 14 | 0.9683264590 | 0.7838296563 | 0.6580566159 | 0.5767359169 |
| 15 | 0.9573907069 | 0.7720219604 | 0.6479480903 | 0.5679433502 |
| 16 | 0.9473481269 | 0.7616242588 | 0.6389600669 | 0.5602246539 |
| 17 | 0.9381934358 | 0.7524384294 | 0.6310986318 | 0.5534476596 |
| 18 | 0.9298621989 | 0.7442777202 | 0.6241512593 | 0.5474899996 |
| 19 | 0.9222036498 | 0.7369202702 | 0.6179423360 | 0.5421397650 |
| 20 | 0.9152117058 | 0.7303213762 | 0.6124022678 | 0.5373821539 |
| 21 | 0.9088144785 | 0.7243886430 | 0.6074011038 | 0.5330849001 |
| 22 | 0.9029047270 | 0.7189831426 | 0.6028683696 | 0.5292093050 |
| 23 | 0.8974529262 | 0.7140279227 | 0.5987629541 | 0.5256806365 |
| 24 | 0.8924307768 | 0.7095202193 | 0.5950086222 | 0.5224711129 |

attempted Richardson extrapolation and other techniques, with no luck as of yet. My hope is that future work will lead to a robust extrapolation approach or other method for approximating the box-counting dimension from the covers $\sigma_{k} \cup \sigma_{k+1}$, and in turn clearly reveal a relationship with the Hausdorff dimension. The combinatorics of the intervals in the canonical covers, discussed in the next section, will aid in this endeavor and expose some theoretically interesting structure.


Figure 3.3 : Plot of $f_{k, \lambda}(\varepsilon)$ for $\lambda=1$ (top) and $\lambda=2$ (bottom). The dashed line indicates the approximate Hausdorff dimension computed with the covers $\sigma_{k} \cup \sigma_{k+1}$, $k=24,25$.


Figure 3.4 : Plot of $f_{k, \lambda}(\varepsilon)$ for $\lambda=3$ (top) and $\lambda=4$ (bottom). The dashed line indicates the approximate Hausdorff dimension computed with the covers $\sigma_{k} \cup \sigma_{k+1}$, $k=24,25$.

### 3.2 Interval Combinatorics

A deep understanding of the combinatorial properties of the bands in $\sigma_{k}$ will provide crucial insight into the spectrum $\Sigma_{\lambda}$. In particular, precise combinatorics will refine extrapolation techniques for fractal dimensions. For example, the work of Damanik et al. demonstrates an application of combinatorial properties for $\lambda>4$ to approximate the box-counting dimension as $\lambda \rightarrow \infty$ [9]. I am seeking statements of the same flavor as Lemma 3.1 below (from [19]) in the intermediate coupling regime $\lambda \in(0,4]$. The authors of [19] first define two types of bands in $\sigma_{k}$.

Definition 3.3. If $I_{k} \subset \sigma_{k}$ is in $\sigma_{k-1}$, then $I_{k}$ is a type $A$ band. If $I_{k} \subset \sigma_{k}$ is in $\sigma_{k-2}$, then $I_{k}$ is a type $B$ band.

Then they proved the following lemma.

Lemma 3.1. Assume $\lambda>4$. Every type $A$ band in $\sigma_{k}$ contains only one type $B$ band in $\sigma_{k+2}$ and no others. Every type $B$ band in $\sigma_{k}$ contains only one type $A$ band in $\sigma_{k+1}$ and two type $B$ bands in $\sigma_{k+2}$ lying on either side of the band in the previous level.

This statement may be visualized in Figure 3.5. It follows immediately that the number of bands in $\sigma_{k} \cup \sigma_{k+1}$ is $2 F_{k}$, provided $\lambda>4$. The proof of Lemma 3.1 relies on the fact that for $\lambda>4$,

$$
\sigma_{k} \cap \sigma_{k+1} \cap \sigma_{k+2}=\emptyset \text { for all } k
$$



Figure 3.5: Bands of $\sigma_{k}$ for $k=1, \ldots, 5$ with $\lambda=5$.


Figure 3.6: Measure of the intersection $\sigma_{k} \cap \sigma_{k+1} \cap \sigma_{k+2}$.
which follows from the trace map and the equivalence $E \in \sigma_{k} \Longleftrightarrow\left|x_{k}(E)\right| \leq 2$ (see Theorem 1.3).

I am interested in the intermediate values of the coupling strength greater than zero but less than 4 ; little is known regarding the combinatorial properties of the bands in $\sigma_{k}$ (and the covers $\sigma_{k} \cup \sigma_{k+1}$ ) in this regime. Can precise statements like Lemma 3.1 be made for smaller values of $\lambda$ ? How do the bands in $\sigma_{k}$ relate to bands in higher levels?

Figure 3.6 verifies a nontrivial intersection of three consecutive $\sigma_{k}$ sets for $\lambda$ values
less than 4. Apparently the measure of the overlap eventually decays exponentially, with more rapid decay for larger $\lambda$ values: the behavior of bands in three consecutive spectra approach the behavior proven for $\lambda$ values greater than 4 . Further, only $\lambda$ values less than 3 are depicted in Figure 3.6 because my calculations suggest that $\sigma_{k} \cap$ $\sigma_{k+1} \cap \sigma_{k+2}=\emptyset$ for $\lambda>3$. How does this overlap behavior affect the combinatorics?

An investigation of the number of intervals in the covers $\sigma_{k} \cup \sigma_{k+1}$ reveals interesting structure. Table 3.4 exhibits the number of intervals in the canonical covers for various $\lambda$ values less than 4 . It appears that the sequence of the number of intervals eventually becomes Fibonacci regardless of the coupling strength, although further calculation is necessary for $\lambda<0.75$. This observation resonates with the intuition that less overlap among consecutive $\sigma_{k}$ sets leads to accumulation of more bands in the covers.

Accurately counting the number of intervals in the numerically obtained set union $\sigma_{k} \cup \sigma_{k+1}$ presents a subtle challenge in light of floating point error and theoretically distinct eigenvalues that appear numerically "close." Hence for completeness, I have included Figure 3.7, which displays the minimum distance between end-points of intervals in $\sigma_{k}$ and end-points of intervals in $\sigma_{k+1}$. The figure indicates several close eigenvalues for smaller $k$ but nothing problematic for larger $k$ that would disrupt the observed Fibonacci recurrence.

In light of the eventual Fibonacci recurrence $\left(N_{k}=N_{k-1}+N_{k-2}\right.$, not necessarily with the convential Fibonacci starting values) depicted in Table 3.4, it is natural to

Table 3.4: Number of intervals in $\sigma_{k} \cup \sigma_{k+1}$.

| $k$ | $\lambda=0.25$ | $\lambda=0.5$ | $\lambda=0.75$ | $\lambda=1$ | $\lambda=1.25$ | $\lambda=1.5$ | $\lambda=1.75$ | $\lambda=2$ | $\lambda=4.1$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 2 |
| 2 | 1 | 1 | 1 | 1 | 2 | 2 | 2 | 2 | 4 |
| 3 | 1 | 1 | 3 | 3 | 3 | 3 | 4 | 4 | 6 |
| 4 | 2 | 3 | 4 | 4 | 5 | 6 | 6 | 6 | 10 |
| 5 | 3 | 5 | 6 | 7 | 8 | 10 | 10 | 10 | 16 |
| 6 | 5 | 7 | 10 | 11 | 13 | 16 | 16 | 16 | 26 |
| 7 | 9 | 13 | 16 | 19 | 22 | 26 | 26 | 26 | 42 |
| 8 | 15 | 21 | 26 | 30 | 36 | 42 | 42 | 42 | 68 |
| 9 | 23 | 33 | 42 | 48 | 58 | 68 | 68 | 68 | 110 |
| 10 | 37 | 53 | 68 | 78 | 94 | 110 | 110 | 110 | 178 |
| 11 | 60 | 87 | 110 | 126 | 152 | 178 | 178 | 178 | 288 |
| 12 | 97 | 141 | 178 | 204 | 246 | 288 | 288 | 288 | 466 |
| 13 | 157 | 227 | 288 | 330 | 398 | 466 | 466 | 466 | 754 |
| 14 | 253 | 367 | 466 | 534 | 644 | 754 | 754 | 754 | 1220 |
| 15 | 411 | 595 | 754 | 864 | 1042 | 1220 | 1220 | 1220 | 1974 |
| 16 | 665 | 963 | 1220 | 1398 | 1686 | 1974 | 1974 | 1974 | 3194 |
| 17 | 1075 | 1557 | 1974 | 2262 | 2728 | 3194 | 3194 | 3194 | 5168 |
| 18 | 1739 | 2520 | 3194 | 3660 | 4414 | 5168 | 5168 | 5168 | 8362 |
| 19 | 2815 | 4078 | 5168 | 5922 | 7142 | 8362 | 8362 | 8362 | 13530 |
| 20 | 4555 | 6598 | 8362 | 9582 | 11556 | 13530 | 13530 | 13530 | 21892 |
| 21 | 7371 | 10676 | 13530 | 15504 | 18698 | 21892 | 21892 | 21892 | 35422 |
| 22 | 11925 | 17274 | 21892 | 25086 | 30254 | 35422 | 35422 | 35422 | 57314 |
| 23 | 19295 | 27950 | 35422 | 40590 | 48952 | 57314 | 57314 | 57314 | 92736 |
| 24 | 31221 | 45224 | 57314 | 65676 | 79206 | 92736 | 92736 | 92736 | 150050 |
| 25 | 50517 | 73174 | 92736 | 106266 | 128158 | 150050 | 150050 | 150050 | 242786 |



Figure 3.7 : Minimum distance between end-points of the intervals in $\sigma_{k}$ and endpoints of the intervals in $\sigma_{k+1}$.


Figure 3.8 : Maximum level $k(\geq 3)$ at which the number of intervals in $\sigma_{k} \cup \sigma_{k+1}$ does not obey the recurrence $N_{k}=N_{k-1}+N_{k-2}$.
investigate the point at which the recurrence appears in the sequence of number of intervals. This is done as follows: for a fixed $\lambda$, build a vector containing the number of intervals in $\sigma_{k} \cup \sigma_{k+1}$ for $k=1, \ldots, 14$. Then, pick the maximum level (starting with $k=3$ ) so that $N_{k} \neq N_{k-1}+N_{k-2}$. Figure 3.8 shows this maximum level for a fine mesh of $\lambda$ values.

One can also investigate the number of intervals in the canonical covers as a function of $\lambda$. The figures given in the survey paper by Damanik, Embree, and Gorodetski elegantly display the canonical covers ( $k$ is fixed) as a functions of $\lambda$ so


Figure 3.9: Number of intervals in $\sigma_{k} \cup \sigma_{k+1}$ for several values of $k$.
that one clearly sees splitting of one interval into multiple intervals as $\lambda$ increases [8]. Figure 3.9 demonstrates that this splitting occurs in a self-similar fashion. It appears the plateaus for a fixed $k$ developing in between the more prominent plateaus form in systematic way, i.e. for $k=8$, see the regions $1<\lambda<1.5$ and $2.5<\lambda<3$.

## Chapter 4

## Concluding Remarks and Future Work

This work presents an $O\left(n^{2}\right)$ algorithm for computing the spectra of periodic Jacobi operators and presents an application to a quasicrystal model, the Fibonacci Hamiltonian. More precisely, the sets $\sigma_{k}$, spectra of related periodic Jacobi operators, form a cover $\sigma_{k} \cup \sigma_{k+1}$ for the Cantor spectrum of the Fibonacci Hamiltonian. Thus, by efficiently computing the sets $\sigma_{k}$ with Algorithm 1, the Cantor spectrum may be approximated with the canonical covers. The sets $\sigma_{k} \cup \sigma_{k+1}$ allow one to numerically grasp the complicated Cantor structure of the spectrum with a goal of informing conjectures and gaining physical insight into quasicrystals.

The application of Algorithm 1 to the Fibonacci model focused on the less-studied intermediate coupling regime $\lambda \in(0,4]$. Future numerical work includes determining a more robust way of extrapolating the box-counting dimension from the canonical covers.

These numerical calculations motivate theoretical questions that require further investigation. What can be said about the regularity and monotonicity of the boxcounting and Hausdorff dimensions in the intermediate coupling regime? Can one develop a closed-form expression for the fractal dimension from plots like Figure 3.1? Further, with perhaps more digits to the approximation in Table 3.2, can one
conjecture and prove at what critical value of $\lambda$ the dimension is equal to $1 / 2$ ? This endeavor naturally leads to further numerical and theoretical investigation of the twoand three-dimensional Fibonacci Hamiltonians.

Future work will also incorporate a theoretical investigation of the combinatorics of the intervals in the sets $\sigma_{k}$ for $\lambda$ in the intermediate coupling regime. I would like to rigorously explain the behavior observed in Figure 3.9 and Table 3.4. Can one prove that for any $\lambda>0$, the number of intervals in the canonical covers eventually obeys the Fibonacci recurrence? As a starting point, one can quickly show that for $\lambda>2$,

$$
\sigma_{k} \cap \sigma_{k+1} \cap \sigma_{k+2} \cap \sigma_{k+3}=\emptyset \text { for all } k
$$

by using the trace map. It would be interesting to use the above result and arguments adapted from Killip et al. to prove a statement similar to Lemma 3.1 [19]. Further, can one show that for $\lambda>3, \sigma_{k} \cap \sigma_{k+1} \cap \sigma_{k+2}=\emptyset$ ?. I hope a better understanding of the combinatorics will help answer questions regarding fractal dimensions and in turn, contribute to our knowledge of quasicrystals.

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[^0]:    David Damanik
    Professor of Mathematics

[^1]:    *I thank Danny Sorensen for showing me the simple proof presented here.

[^2]:    ${ }^{\dagger}$ Fortran code and subroutines may be compiled to run in quadruple precision by promoting all double precision data types to real*16. This was accomplished in my implementation with the following script: http://icl.cs.utk.edu/lapack-forum/viewtopic.php?f=2\&t=2739
    ${ }^{\ddagger}$ In my testing of these matrices from Gear, I used an sorting subroutine from the following website: http://jean-pierre.moreau.pagesperso-orange.fr/

